

CS388: Natural Language Processing

Lecture 14: Semantics I

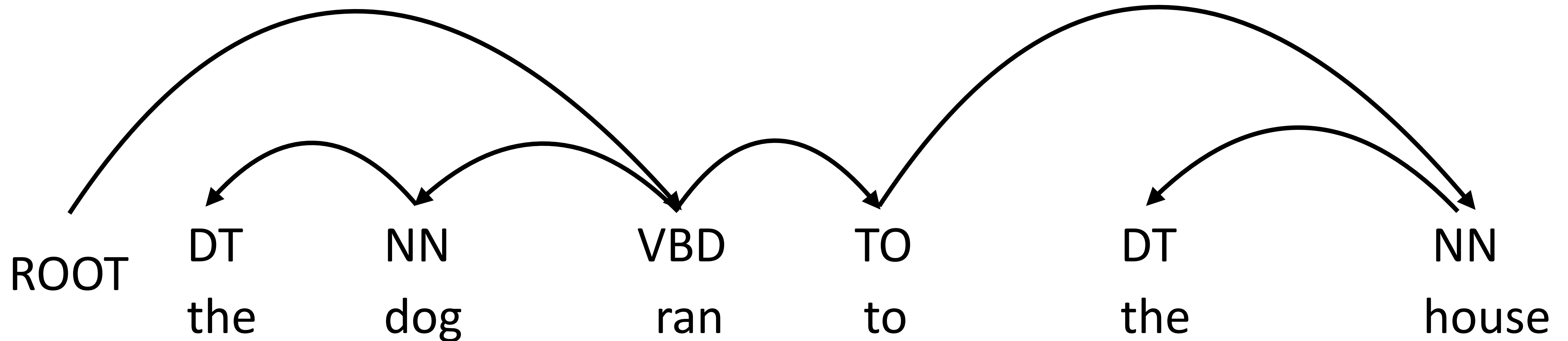
Greg Durrett





Recall: Dependencies

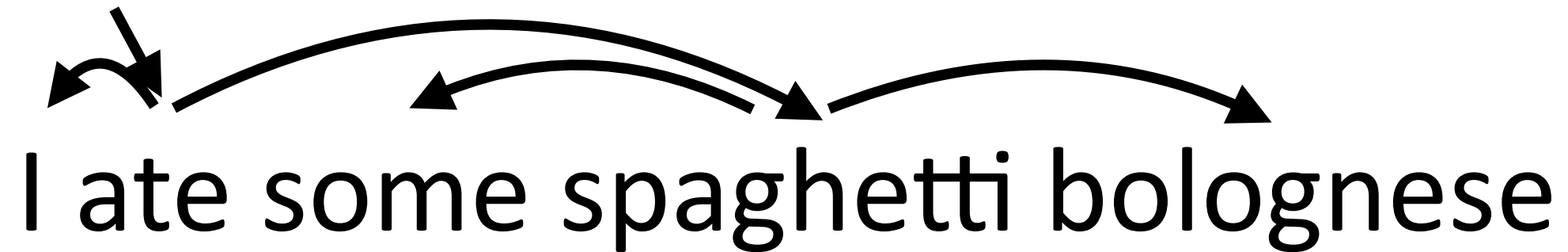
- ▶ Dependency syntax: syntactic structure is defined by dependencies
 - ▶ Head (parent, governor) connected to dependent (child, modifier)
 - ▶ Each word has exactly one parent except for the ROOT symbol
 - ▶ Dependencies must form a directed acyclic graph





Recall: Shift-Reduce Parsing

ROOT



► State: **Stack:** [ROOT I ate] **Buffer:** [some spaghetti bolognese]

► Left-arc (reduce operation): Let σ denote the stack

► “Pop two elements, add an arc, put them back on the stack”

$$\boxed{\sigma | w_{-2}, w_{-1}} \rightarrow \boxed{\sigma | w_{-1}}, \quad w_{-2} \text{ is now a child of } w_{-1}$$

► Train a classifier to make these decisions sequentially — that classifier can parse sentences for you



Where are we now?

- ▶ Early in the class: bags of word (classifiers) => sequences of words (sequence modeling)
- ▶ Now we can understand sentences in terms of tree structures as well
- ▶ Why is this useful? What does this allow us to do?
- ▶ We're going to see how parsing can be a stepping stone towards more formal representations of language meaning



Today

- ▶ Montague semantics:
 - ▶ Model theoretic semantics
 - ▶ Compositional semantics with first-order logic
- ▶ CCG parsing for database queries
- ▶ Lambda-DCS for question answering

Model Theoretic Semantics



Model Theoretic Semantics

- ▶ Key idea: can ground out natural language expressions in set-theoretic expressions called *models* of those sentences
- ▶ Natural language statement $S \Rightarrow$ interpretation of S that models it
She likes going to that restaurant
 - ▶ Interpretation: defines who *she* and *that restaurant* are, make it able to be concretely evaluated with respect to a *world*
- ▶ Entailment (statement A implies statement B) reduces to: in all worlds where A is true, B is true
- ▶ Our modeling language is *first-order logic*



First-order Logic

- ▶ Powerful logic formalism including things like entities, relations, and quantifications

Lady Gaga sings

- ▶ sings is a *predicate* (with one argument), function $f: \text{entity} \rightarrow \text{true/false}$
- ▶ sings(Lady Gaga) = true or false, have to execute this against some database (*world*)
- ▶ $[[\text{sings}]] = \text{denotation}$, set of entities which sing (found by executing this predicate on the world — we'll come back to this)



Quantification

- ▶ Universal quantification: “forall” operator

- ▶ $\forall x \text{ sings}(x) \vee \text{dances}(x) \rightarrow \text{performs}(x)$

“Everyone who sings or dances performs”

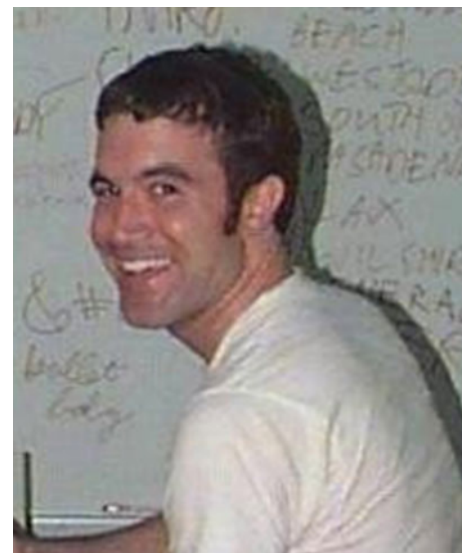
- ▶ Existential quantification: “there exists” operator

- ▶ $\exists x \text{ sings}(x)$ *“Someone sings”*

- ▶ Source of ambiguity! *“Everyone is friends with someone”*

- ▶ $\forall x \exists y \text{ friend}(x,y)$

- ▶ $\exists y \forall x \text{ friend}(x,y)$





Logic in NLP

- ▶ Question answering:

Who are all the American singers named Amy?

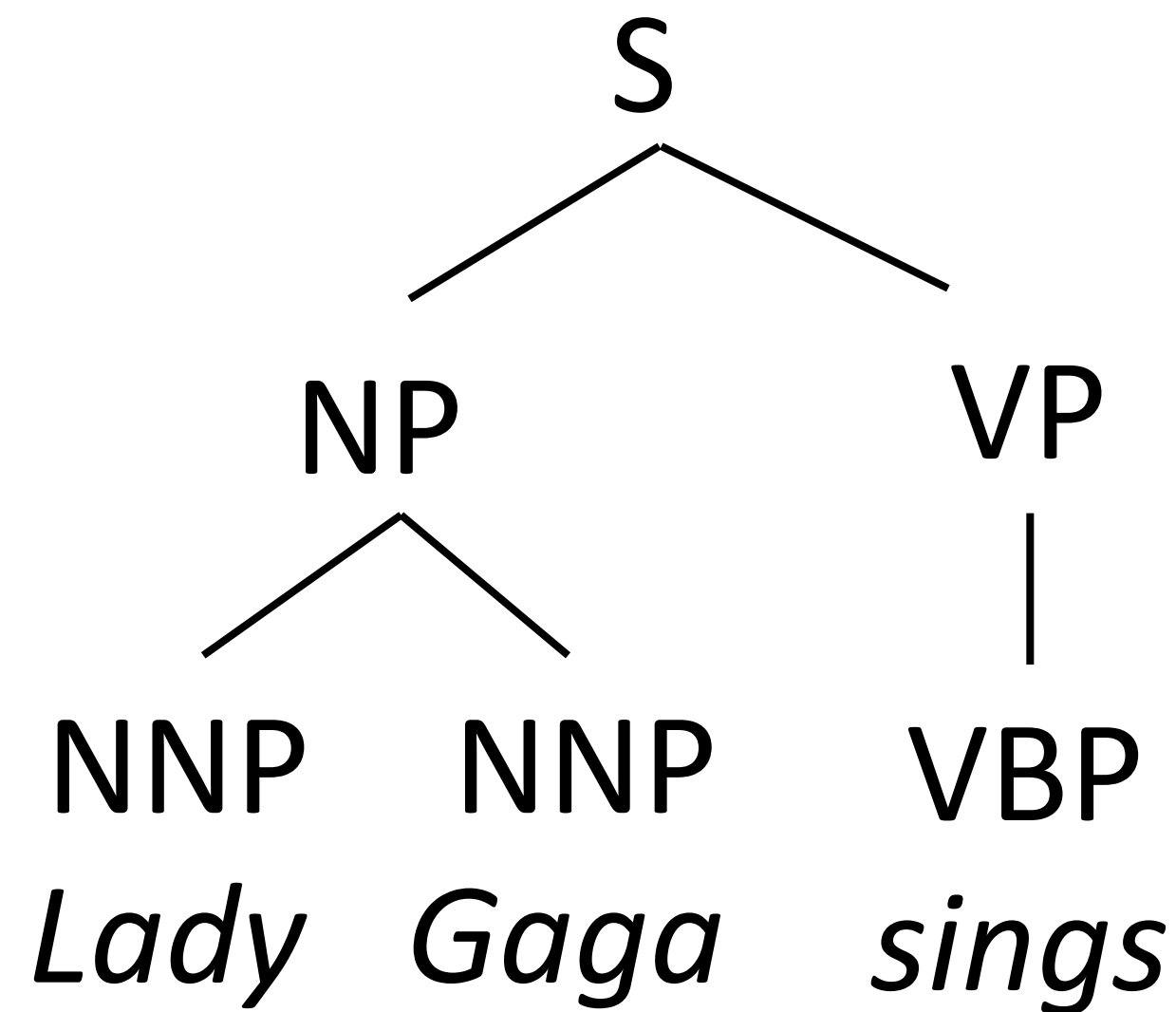
$\lambda x. \text{nationality}(x, \text{USA}) \wedge \text{sings}(x) \wedge \text{firstName}(x, \text{Amy})$

- ▶ Function that maps from x to true/false, like `filter`. Execute this on the world to answer the question
- ▶ Lambda calculus: powerful system for expressing these functions
- ▶ Information extraction: *Lady Gaga and Eminem are both musicians*
 $\text{musician}(\text{Lady Gaga}) \wedge \text{musician}(\text{Eminem})$
- ▶ Can now do reasoning. Maybe know: $\forall x \text{ musician}(x) \Rightarrow \text{performer}(x)$
Then: $\text{performer}(\text{Lady Gaga}) \wedge \text{performer}(\text{Eminem})$

Compositional Semantics with First- Order Logic



Montague Semantics



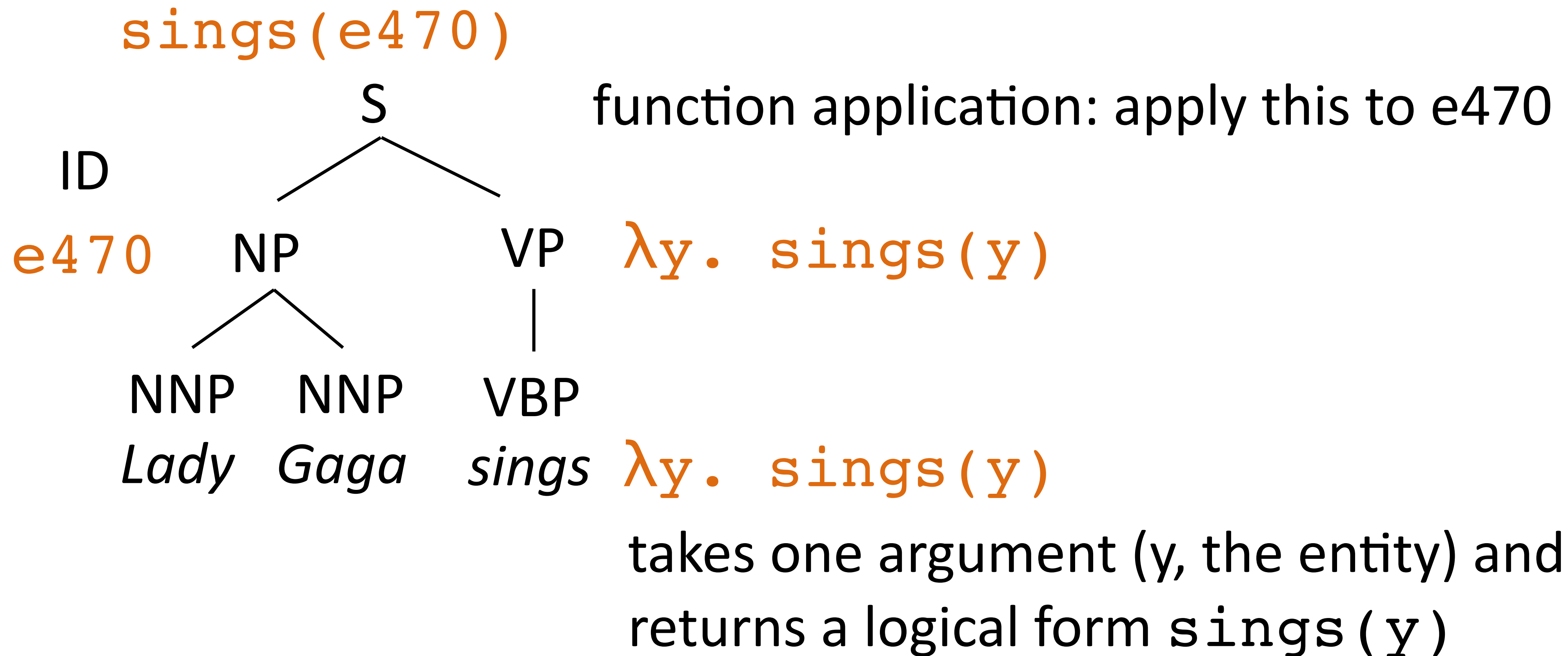
Id	Name	Alias	Birthdate	Sings?
e470	Stefani Germanotta	Lady Gaga	3/28/1986	T
e728	Marshall Mathers	Eminem	10/17/1972	T

► Database containing entities, predicates, etc.

- Sentence expresses something about the world which is either true or false
- Denotation: evaluation of some expression against this database
 - $[[\textit{Lady Gaga}]] = e470$
denotation of this string is an entity
 - $[[\textit{sings}(e470)]] = \text{True}$
denotation of this expression is T/F



Montague Semantics

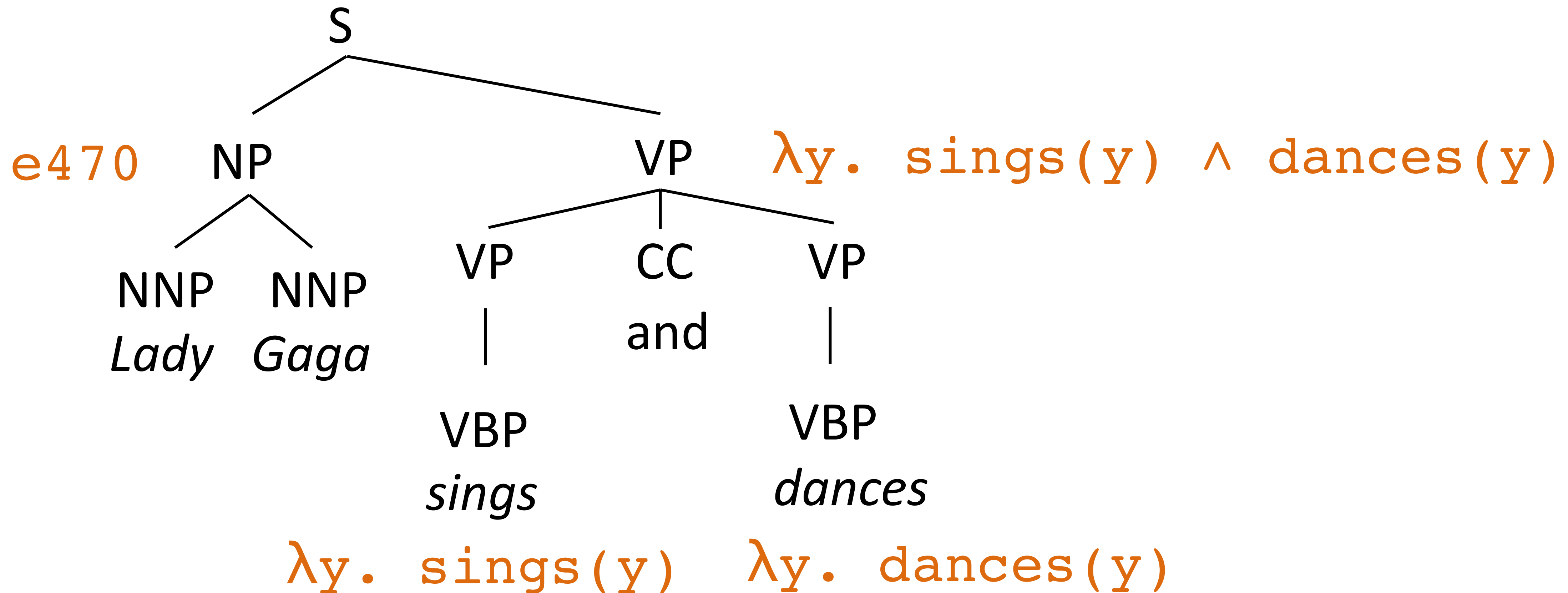


- We can use the syntactic parse as a bridge to the lambda-calculus representation, build up a logical form (our model) *compositionally*



Parses to Logical Forms

$\text{sings}(\text{e470}) \wedge \text{dances}(\text{e470})$

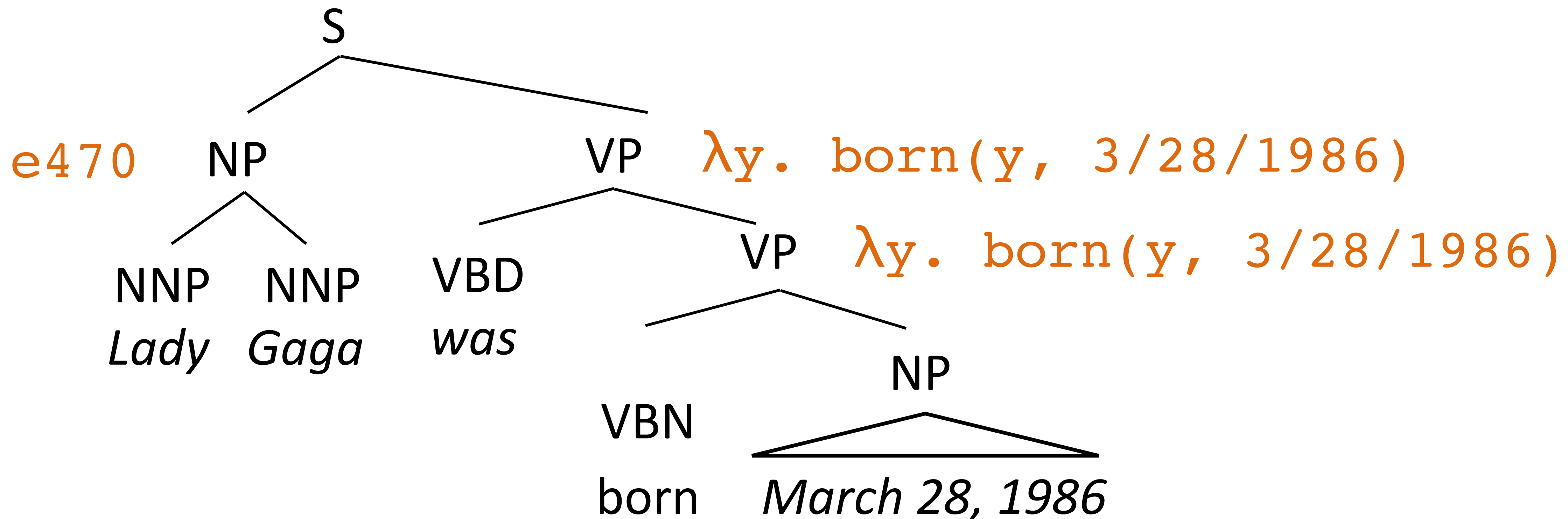


- General rules:
- VP: $\lambda y. a(y) \wedge b(y) \rightarrow \text{VP: } \lambda y. a(y) \text{ CC VP: } \lambda y. b(y)$
 - S: $f(x) \rightarrow \text{NP: } x \text{ VP: } f$



Parses to Logical Forms

$\text{born}(e470, 3/28/1986)$



$\lambda x. \lambda y. \text{born}(y, x) \quad 3/28/1986$

- Function takes two arguments: first x (date), then y (entity)
- How to handle tense: should we indicate that this happened in the past?



Tricky things

- ▶ Adverbs/temporality: *Lady Gaga sang well yesterday*

`sings(Lady Gaga, time=yesterday, manner=well)`

- ▶ “Neo-Davidsonian” view of events: things with many properties:

`∃e. type(e, sing) ∧ agent(e, e470) ∧ manner(e, well) ∧ time(e, ...)`

- ▶ Quantification: *Everyone is friends with someone*

<code>∃y ∀x friend(x, y)</code>	<code>∀x ∃y friend(x, y)</code>
(one friend)	(different friends)

- ▶ Same syntactic parse for both! So syntax doesn't resolve all ambiguities

- ▶ Indefinite: *Amy ate a waffle* `∃w. waffle(w) ∧ ate(Amy, w)`

- ▶ Generic: *Cats eat mice* (all cats eat mice? most cats? some cats?)



Semantic Parsing

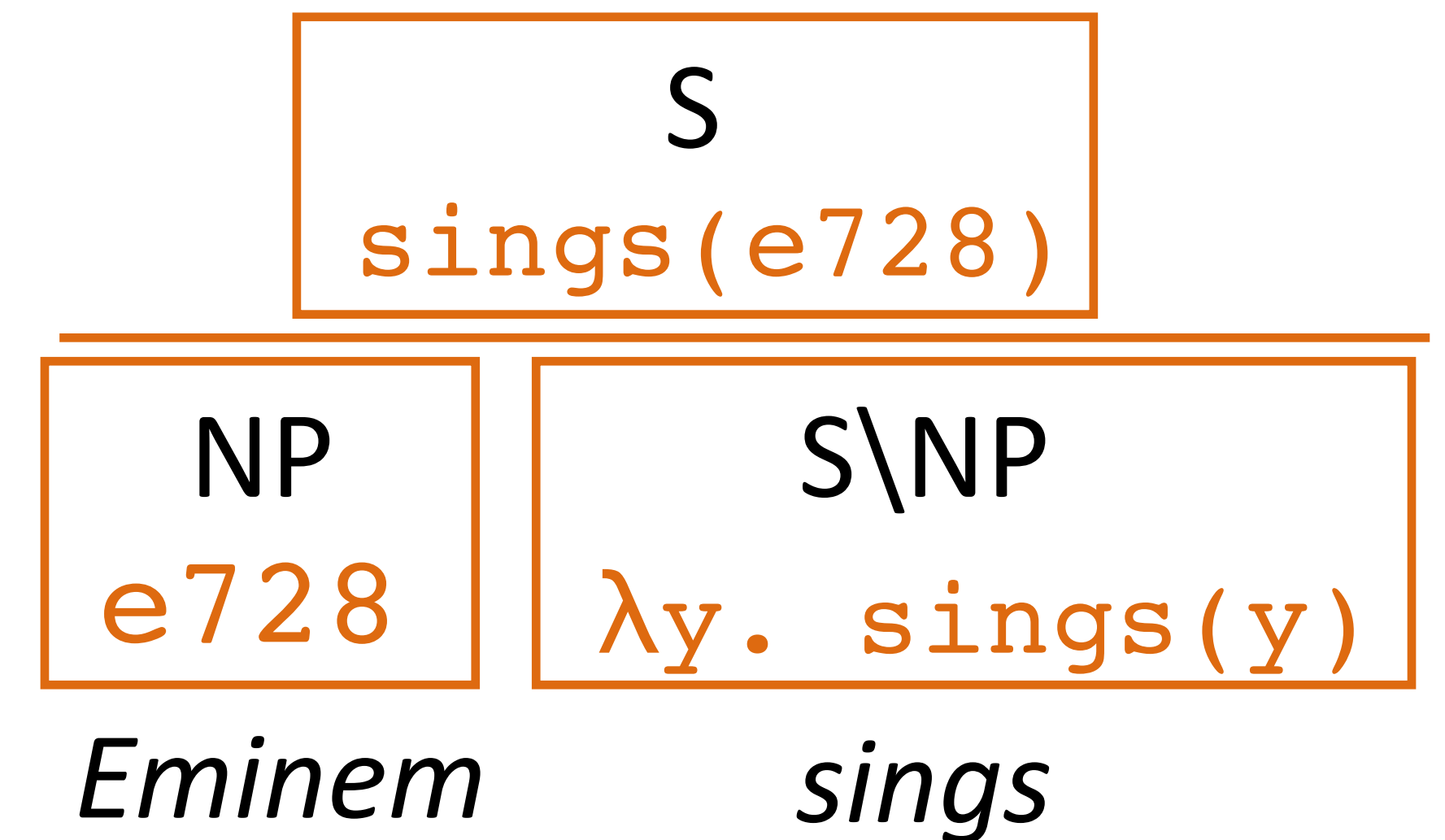
- ▶ For question answering, syntactic parsing doesn't tell you everything you want to know, but indicates the right structure
- ▶ Solution: *semantic parsing*: many forms of this task depending on semantic formalisms
- ▶ Two today: CCG (looks like what we've been doing) and lambda-DCS
- ▶ Applications: database querying/question answer: produce lambda-calculus expressions that can be executed in these contexts

CCG Parsing



Combinatory Categorical Grammar

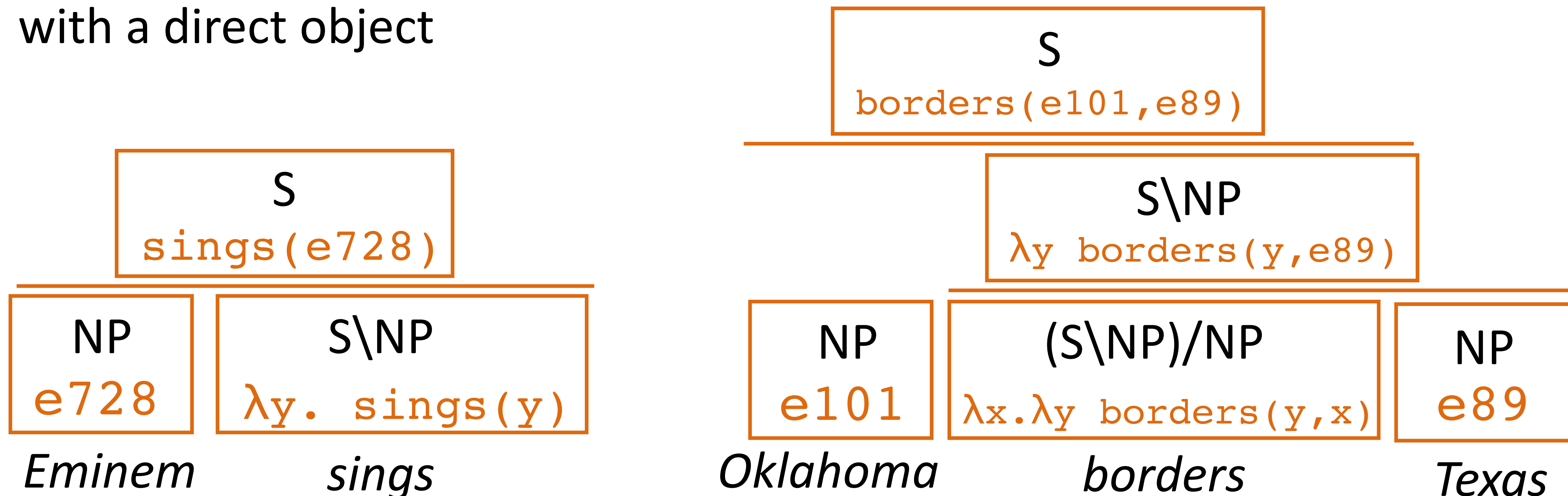
- ▶ Steedman+Szabolcsi (1980s): formalism bridging syntax and semantics
- ▶ Parallel derivations of syntactic parse and lambda calculus expression
- ▶ Syntactic categories (for this lecture): S, NP, “slash” categories
- ▶ $S \backslash NP$: “if I combine with an NP on my left side, I form a sentence” — verb
- ▶ When you apply this, there has to be a parallel instance of function application on the semantics side





Combinatory Categorical Grammar

- ▶ Steedman+Szabolcsi 1980s: formalism bridging syntax and semantics
- ▶ Syntactic categories (for this lecture): S, NP, “slash” categories
 - ▶ $S \backslash NP$: “if I combine with an NP on my left side, I form a sentence” — verb
 - ▶ $(S \backslash NP) / NP$: “I need an NP on my right and then on my left” — verb with a direct object





CCG Parsing

What	states	border	Texas
$(S/(S \backslash NP))/N$ $\lambda f. \lambda g. \lambda x. f(x) \wedge g(x)$	N $\lambda x. state(x)$	$(S \backslash NP)/NP$ $\lambda x. \lambda y. borders(y, x)$	NP $texas$
			$\xrightarrow{\hspace{10em}}$
			$(S \backslash NP)$ $\lambda y. borders(y, texas)$

- ▶ “What” is a **very** complex type: needs a noun and needs a $S \backslash NP$ to form a sentence. $S \backslash NP$ is basically a verb phrase (*border Texas*)



CCG Parsing

What	states	border	Texas
$(S/(S \setminus NP))/N$ $\lambda f. \lambda g. \lambda x. f(x) \wedge g(x)$	N $\lambda x. state(x)$	$(S \setminus NP)/NP$ $\lambda x. \lambda y. borders(y, x)$	NP $texas$
$S/(S \setminus NP)$ $\lambda g. \lambda x. state(x) \wedge g(x)$		$(S \setminus NP)$ $\lambda y. borders(y, texas)$	
S $\lambda x. state(x) \wedge borders(x, texas)$			

- ▶ “What” is a **very** complex type: needs a noun and needs a $S \setminus NP$ to form a sentence. $S \setminus NP$ is basically a verb phrase (*border Texas*)
 - ▶ Lexicon is highly ambiguous — all the challenge of CCG parsing is in picking the right lexicon entries
- Zettlemoyer and Collins (2005)



CCG Parsing

Show me	flights	to	Prague
S/N $\lambda f.f$	N $\lambda x.flight(x)$	$(N \backslash N) / NP$ $\lambda y.\lambda f.\lambda x.f(y) \wedge to(x, y)$	NP PRG
		$N \backslash N$ $\lambda f.\lambda x.f(x) \wedge to(x, PRG)$	
		N $\lambda x.flight(x) \wedge to(x, PRG)$	
		S $\lambda x.flight(x) \wedge to(x, PRG)$	

- “to” needs an NP (destination) and N (parent)



CCG Parsing

- ▶ Many ways to build these parsers
- ▶ One approach: run a “supertagger” (tags the sentence with complex labels), then run the parser

What	states	border	Texas
$\frac{(S/(S \backslash NP))/N}{\lambda f. \lambda g. \lambda x. f(x) \wedge g(x)}$	$\frac{N}{\lambda x. state(x)}$	$\frac{(S \backslash NP)/NP}{\lambda x. \lambda y. borders(y, x)}$	$\frac{NP}{texas}$

- ▶ Parsing is easy once you have the tags, so we’ve reduced it to a (hard) tagging problem

Zettlemoyer and Collins (2005)



Building CCG Parsers

- Model: log-linear model over derivations with features on rules:

$$P(d|x) \propto \exp w^\top \left(\sum_{r \in d} f(r, x) \right)$$

$$\begin{array}{c} f \left(\begin{array}{c} \boxed{\begin{array}{c} S \\ sings(e728) \end{array}} \end{array} \right) = \text{Indicator}(S \rightarrow NP \ S \backslash NP) \\ \hline f \left(\begin{array}{c} \boxed{\begin{array}{c} NP \\ e728 \end{array}} \end{array} \right) \quad f \left(\begin{array}{c} \boxed{\begin{array}{c} S \backslash NP \\ \lambda y. sings(y) \end{array}} \end{array} \right) = \text{Indicator}(S \backslash NP \rightarrow sings) \\ \textit{Eminem} \qquad \qquad \textit{sings} \end{array}$$

- Can parse with a variant of CKY

Zettlemoyer and Collins (2005)



Building CCG Parsers

- ▶ Training data looks like pairs of sentences and logical forms

What states border Texas $\lambda x. \text{state}(x) \wedge \text{borders}(x, \text{e89})$

- ▶ Problem: we don't know the derivation
 - ▶ *Texas* corresponds to NP | **e89** in the logical form (easy to figure out)
 - ▶ *What* corresponds to (S/(S\NP))/N | **$\lambda f. \lambda g. \lambda x. f(x) \wedge g(x)$**
 - ▶ How do we infer that without being told it?



Lexicon

- ▶ GENLEX: takes sentence S and logical form L . Break up logical form into chunks $C(L)$, assume any substring of S might map to any chunk

What states border Texas $\lambda x. \text{state}(x) \wedge \text{borders}(x, e89)$

- ▶ Chunks inferred from the logic form based on rules:
 - ▶ NP: $e89$ ▶ $(S \backslash NP) / NP: \lambda x. \lambda y. \text{borders}(x, y)$
- ▶ Any substring can parse to any of these in the lexicon
 - ▶ *Texas* \rightarrow NP: $e89$ is correct
 - ▶ *border Texas* \rightarrow NP: $e89$
 - ▶ *What states border Texas* \rightarrow NP: $e89$

...

Zettlemoyer and Collins (2005)



GENLEX

Rules		Categories produced from logical form
Input Trigger	Output Category	$\arg \max(\lambda x.state(x) \wedge borders(x, texas), \lambda x.size(x))$
constant c	$NP : c$	$NP : texas$
arity one predicate p_1	$N : \lambda x.p_1(x)$	$N : \lambda x.state(x)$
arity one predicate p_1	$S \backslash NP : \lambda x.p_1(x)$	$S \backslash NP : \lambda x.state(x)$
arity two predicate p_2	$(S \backslash NP) / NP : \lambda x.\lambda y.p_2(y, x)$	$(S \backslash NP) / NP : \lambda x.\lambda y.borders(y, x)$
arity two predicate p_2	$(S \backslash NP) / NP : \lambda x.\lambda y.p_2(x, y)$	$(S \backslash NP) / NP : \lambda x.\lambda y.borders(x, y)$
arity one predicate p_1	$N / N : \lambda g.\lambda x.p_1(x) \wedge g(x)$	$N / N : \lambda g.\lambda x.state(x) \wedge g(x)$
literal with arity two predicate p_2 and constant second argument c	$N / N : \lambda g.\lambda x.p_2(x, c) \wedge g(x)$	$N / N : \lambda g.\lambda x.borders(x, texas) \wedge g(x)$
arity two predicate p_2	$(N \backslash N) / NP : \lambda x.\lambda g.\lambda y.p_2(x, y) \wedge g(x)$	$(N \backslash N) / NP : \lambda g.\lambda x.\lambda y.borders(x, y) \wedge g(x)$
an $\arg \max$ / \min with second argument arity one function f	$NP / N : \lambda g.\arg \max / \min(g, \lambda x.f(x))$	$NP / N : \lambda g.\arg \max(g, \lambda x.size(x))$
an arity one numeric-ranged function f	$S / NP : \lambda x.f(x)$	$S / NP : \lambda x.size(x)$

- Very complex and hand-engineered way of taking lambda calculus expressions and “backsolving” for the derivation

Zettlemoyer and Collins (2005)



Learning

- ▶ Iterative procedure like the EM algorithm: estimate “best” parses that derive each logical form, retrain the parser using these parses with supervised learning
- ▶ We’ll talk about a simpler form of this in a few slides



Applications

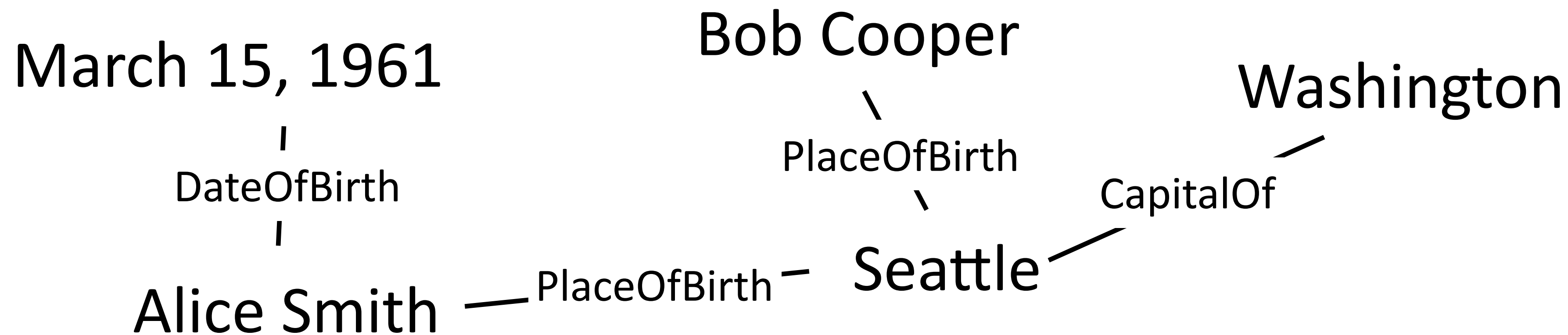
- ▶ GeoQuery: answering questions about states (~80% accuracy)
- ▶ Jobs: answering questions about job postings (~80% accuracy)
- ▶ ATIS: flight search
- ▶ Can do well on all of these tasks if you handcraft systems and use plenty of training data: these domains aren't that rich
- ▶ What about broader QA?

Lambda-DCS



Lambda-DCS

- ▶ Dependency-based compositional semantics — original version was less powerful than lambda calculus, lambda-DCS is as powerful
- ▶ Designed in the context of building a QA system from Freebase
- ▶ Freebase: set of entities and relations



- ▶ $[[\text{PlaceOfBirth}]]$ = set of pairs of (person, location)



Lambda-DCS

Lambda-DCS

Seattle

PlaceOfBirth

PlaceOfBirth.Seattle

Lambda calculus

$\lambda x. x = \text{Seattle}$

$\lambda x. \lambda y. \text{PlaceOfBirth}(x, y)$

$\lambda x. \text{PlaceOfBirth}(x, \text{Seattle})$

- Looks like a tree fragment over Freebase

??? — PlaceOfBirth — Seattle

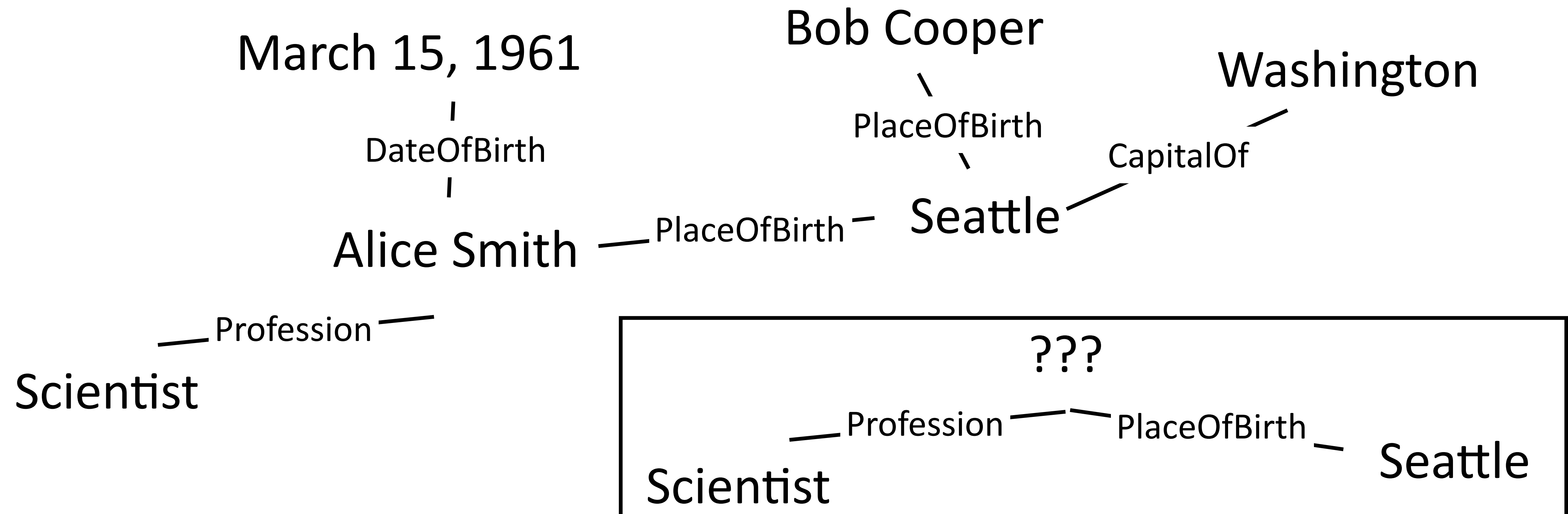
Profession.Scientist \wedge
PlaceOfBirth.Seattle

$\lambda x. \text{Profession}(x, \text{Scientist})$
 $\wedge \text{PlaceOfBirth}(x, \text{Seattle})$

Liang et al. (2011), Liang (2013)



Lambda-DCS



“list of scientists born in Seattle”

`Profession.Scientist ^ PlaceOfBirth.Seattle`

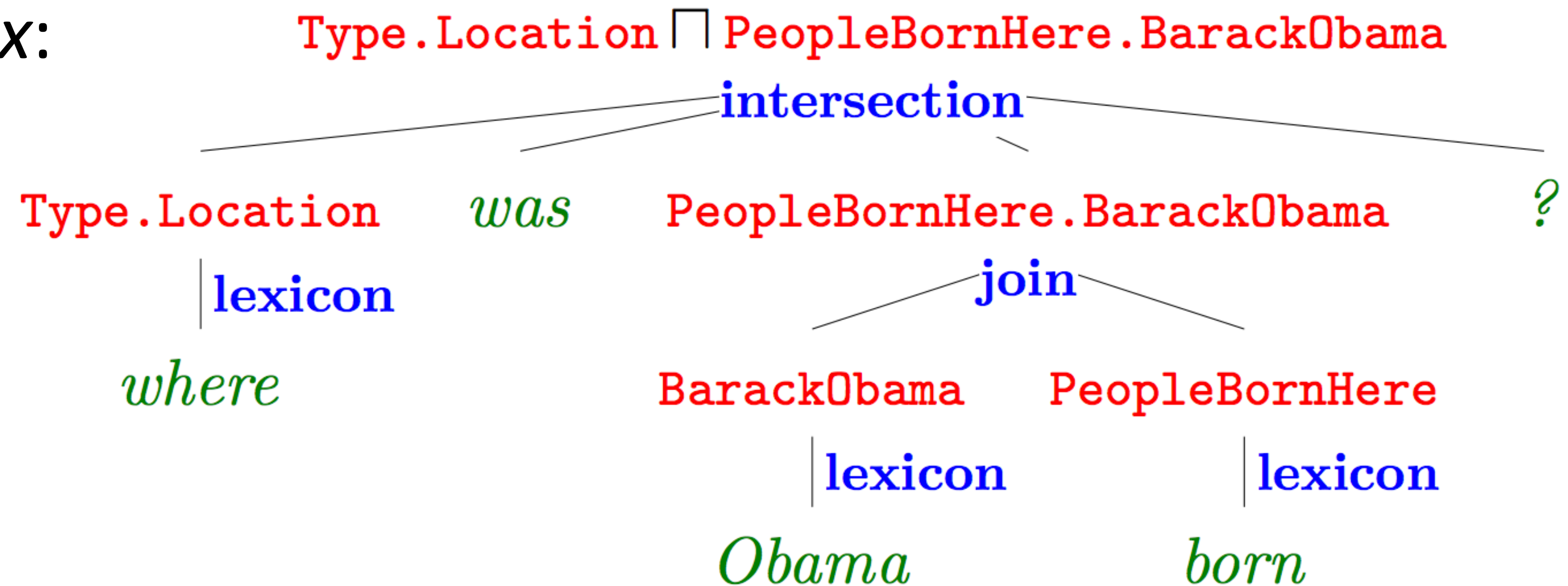
- Execute this fragment against Freebase, returns Alice Smith (and others)

Liang et al. (2011), Liang (2013)



Parsing into Lambda-DCS

- Derivation d on sentence x :



- Building the lexicon: more sophisticated process than GENLEX, but can handle thousands of predicates
- Log-linear model with features on rules:
$$P(d|x) \propto \exp w^\top \left(\sum_{r \in d} f(r, x) \right)$$
- Similar to CRF parsers



Parsing with Lambda-DCS

- ▶ Learn just from question-answer pairs: maximize the likelihood of the right denotation y with the derivation d marginalized out

$$\mathcal{O}(\theta) = \sum_{i=1}^n \log \sum_{d \in D(x) : \llbracket d.z \rrbracket_K = y_i} p_{\theta}(d \mid x_i).$$

sum over derivations d such that the denotation of d on knowledge base K is y_i

For each example:

Run beam search to get a set of derivations

Let d = highest-scoring derivation in the beam

Let d^* = highest-scoring derivation in the beam *with correct denotation*

Do a structured perceptron update towards d^* away from d



Takeaways

- ▶ Can represent meaning with first order logic and lambda calculus
- ▶ Can bridge syntax and semantics and create semantic parsers that can interpret language into lambda-calculus expressions
- ▶ Useful for querying databases, question answering, etc.
- ▶ Next time: neural net methods for doing this that rely less on having explicit grammars